

Markov-chain Monte Carlo: A modern primer 2/2

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W. Krauth; Oxford University Press (2006)
Statistical mechanics: Algorithms and computations

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- 1 Convergence theorem—A priori probabilities
- 2 Perfect sampling—coupling
- 3 (Meta algorithms—extended ensembles)

Markov-chain convergence theorem

For P irreducible and aperiodic, with stationary distribution π :

$$\max_{x \in \Omega} \|P(x, \cdot) - \pi\|_{\text{TV}} \leq C \alpha^t$$

with $C > 0$ and $\alpha \in (0, 1)$.

- Exponential convergence is everywhere, but C and α are unknown.
- Can we do better?

Converging faster than exponential

- 1 Absorbing Markov chain with one absorbing state.

$$P = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \end{pmatrix}$$

- 2 (Starting with $\pi^{\{0\}} = \pi$.)
- 3 Transition matrix $P_{ij} = \pi_j$.

$$\pi_i^{\{t+1\}} = \sum_j \pi_j^{\{t\}} P_{ji} = \underbrace{\sum_j \pi_j^{\{t\}}}_{=1} \pi_i$$

Convergence in one step, better than exponential.

Metropolis–Hastings algorithm (1/2)

$$P(a \rightarrow b) = \underbrace{\mathcal{A}(a \rightarrow b)}_{\text{consider } a \rightarrow b} \cdot \underbrace{\mathcal{P}(a \rightarrow b)}_{\text{accept } a \rightarrow b}.$$

Detailed balance:

$$\pi(a)P(a \rightarrow b) = \pi(b)P(b \rightarrow a) \quad (1)$$

$$\frac{\mathcal{P}(a \rightarrow b)}{\mathcal{P}(b \rightarrow a)} = \frac{\pi(b)}{\mathcal{A}(a \rightarrow b)} \frac{\mathcal{A}(b \rightarrow a)}{\pi(a)}.$$

This leads to a generalized Metropolis filter

$$\mathcal{P}(a \rightarrow b) = \min \left[1, \frac{\pi(b)}{\mathcal{A}(a \rightarrow b)} \frac{\mathcal{A}(b \rightarrow a)}{\pi(a)} \right]$$

Metropolis–Hastings algorithm (2/2)

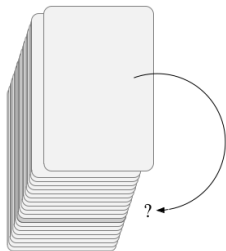
- Generalized Metropolis filter

$$\mathcal{P}(a \rightarrow b) = \min \left[1, \frac{\pi(b)}{\mathcal{A}(a \rightarrow b)} \frac{\mathcal{A}(b \rightarrow a)}{\pi(a)} \right]$$

- $\mathcal{A}(a \rightarrow b) = \pi(b)$ unrealistic
- $\mathcal{A}(a \rightarrow b) \simeq \pi(b)$ realistic, super interesting.
- MCMC equivalent of perturbation theory in theoretical physics.
- Better \mathcal{A} 's \Leftrightarrow larger moves.
- Applications in spin models, bosonic QMC, etc.

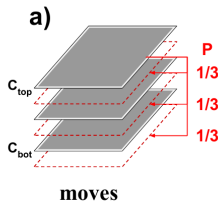
Identify good \mathcal{A} 's through machine learning?

Shuffling of cards 1/5



- $\Omega_n^{\text{shuffle}} = \{\text{Permutations of } \{1, \dots, n\}\}$
- For $n = 3$:
 $\Omega_3^{\text{shuffle}} = \{1 \equiv \{1, 2, 3\}, 2 \equiv \{1, 3, 2\}, 3 \equiv \{2, 1, 3\}, 4 \equiv \{2, 3, 1\}, 5 \equiv \{3, 1, 2\}, 6 \equiv \{3, 2, 1\}\}.$
- $\pi^{t=0} = \delta(\{1, \dots, n\})$ (perfectly ordered set)

Shuffling of cards 2/5

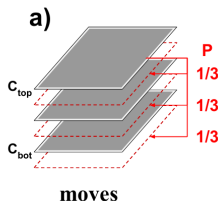


- $\Omega_3^{\text{shuffle}} = \{1 \equiv \{1, 2, 3\}, 2 \equiv \{1, 3, 2\}, 3 \equiv \{2, 1, 3\}, 4 \equiv \{2, 3, 1\}, 5 \equiv \{3, 1, 2\}, 6 \equiv \{3, 2, 1\}\}.$

•

$$P = \frac{1}{3} \begin{pmatrix} 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 & 1 \end{pmatrix}$$

Shuffling of cards 3/5



procedure top-to-random

input $\{c_1, \dots, c_n\}$

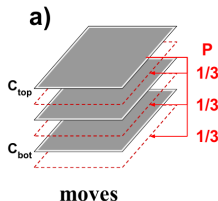
$i \leftarrow \text{choice}(\{1, \dots, n\})$

$\{\hat{c}_1, \dots, \hat{c}_n\} \leftarrow \{c_2, \dots, c_i, c_1, c_{i+1}, \dots, c_n\}$

output $\{\hat{c}_1, \dots, \hat{c}_n\}$

- Insert upper card (c_1) behind card i and before card $i + 1$
- NB: if $i = 1$, put it back on top.

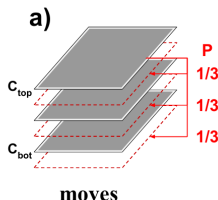
Shuffling of cards 4/5



```
procedure top2random-stop
input  $\{c_1, \dots, c_n\}$ 
 $c_{\text{first-}n} \leftarrow c_n$ 
for  $t = 1, 2, \dots$  do
     $\tilde{c}_1 \leftarrow c_1$ 
     $\{c_1, \dots, c_n\} \leftarrow \text{top2random}(\{c_1, \dots, c_n\})$ 
    if  $(\tilde{c}_1 = c_{\text{first-}n})$  break
output  $\{c_1, \dots, c_n, t\}$ 
```

- Perfect sample (!).
- Expected running time: $n \log n$.

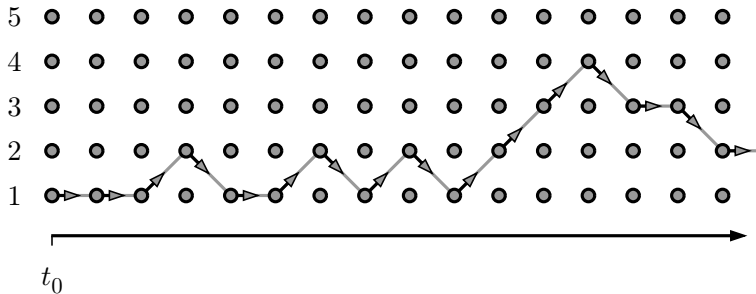
Shuffling of cards 5/5



```
procedure direct-shuffle
input  $\{c_1, \dots, c_n\}$ 
for  $t = 1, \dots, n$  do
     $i \leftarrow \text{choice}(\{n - t + 1, \dots, n\})$ 
     $\{c_1, \dots, c_n\} \leftarrow \{c_2, \dots, c_i, c_1, c_{i+1}, \dots, c_n\}$ 
output  $\{c_1, \dots, c_n\}$ 
```

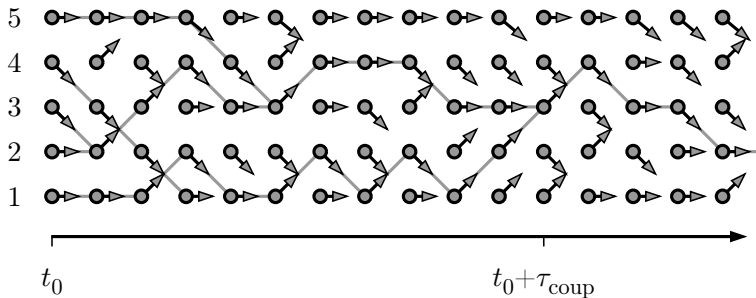
- Running time: n .
- Running time: n .
- Standard algorithm for generating random permutations.

Markov chain (traditional view)



- Configuration c_t , move δ_t .
- Set $t_0 = 0$.

Markov chain (random maps), coupling 1/4



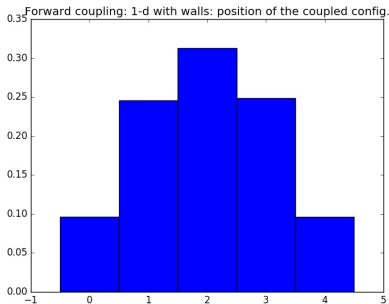
- Each configuration has its move at each time step.
- Coupling (Doeblin, 1930s).

Markov chain (random maps), coupling 2/4

```
pos=[]
for stat in range(10000):
    posit=set(range(N))
    for t in range(1000000):
        posit = set([min(max(b + random.randint(-1, 1), 0), N - 1) for b in posit])
        if len(posit) == 1: break
    pos.append(posit.pop())
```

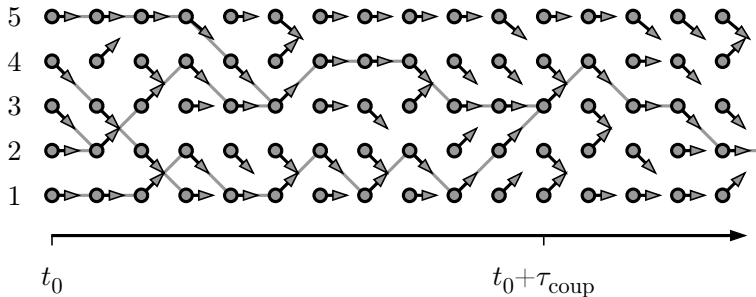
- Position of coupling not uniform.
- Coupling time larger than mixing time.

Markov chain (random maps), coupling 3/4



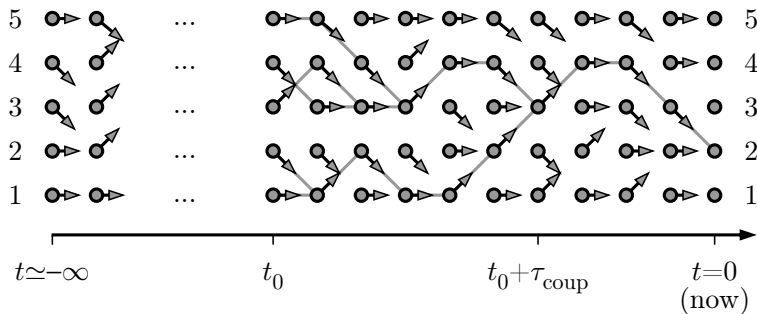
- Histogram of coupling position.

Markov chain (random maps), coupling 4/4



- Each configuration has its move at each time step.
- Coupling (Doeblin, 1930s).

Coupling from the past 1/8



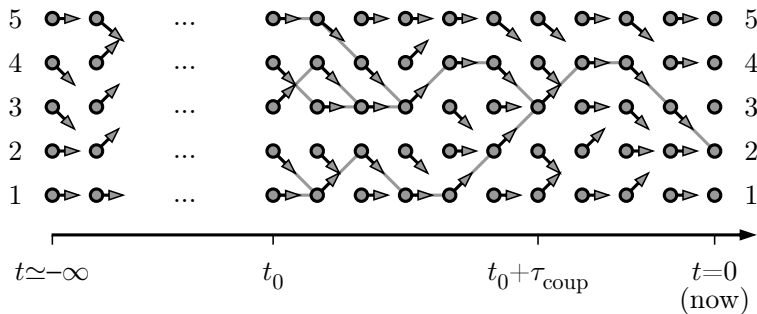
- Starting an MCMC simulation at $t = -\infty$
- Propp & Wilson (1997)

Coupling from the past 2/8

```
pos = []
for statistic in range(1000000):
    all_arrows = {}
    time_tot = 0
    while True:
        time_tot -= 1
        arrows = [random.randint(-1, 1) for i in range(N)]
        if arrows[0] == -1: arrows[0] = 0
        if arrows[N - 1] == 1: arrows[N - 1] = 0
        all_arrows[time_tot] = arrows
        positions = set(range(0, N))
        for t in range(time_tot, 0):
            positions = set([b + all_arrows[t][b] for b in positions])
            if len(positions) == 1: break
        if len(positions) == 1: break
```

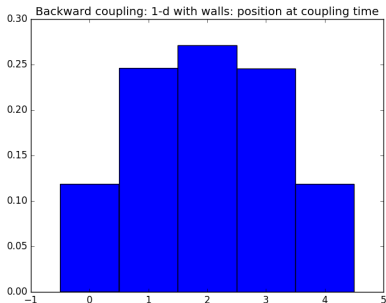
- Starting an MCMC simulation at $t = -\infty$
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Coupling from the past 3/8



- Starting an MCMC simulation at $t = -\infty$
- Propp & Wilson (1997)

Coupling from the past 4/8



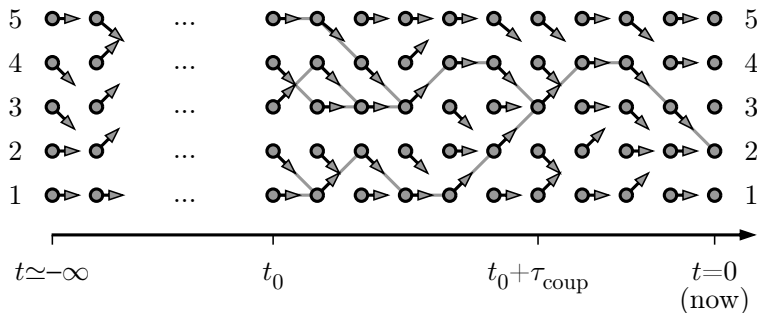
- Coupling position (in the past) non-uniform)

Coupling from the past 5/8

```
for statistic in range(10000):
    all_arrows = {}
    time_tot = 0
    while True:
        time_tot -= 1
        old_pos = set(range(0, N))
        arrows = [random.randint(-1, 1) for i in range(N)]
        if arrows[0] == -1: arrows[0] = 0
        if arrows[N - 1] == 1: arrows[N - 1] = 0
        all_arrows[time_tot] = arrows
        positions = set(range(N))
        for t in range(time_tot, 0):
            positions = set([b + all_arrows[t][b] for b in positions])
        if len(positions) == 1: break
    a=positions.pop()
    pos.append(a)
```

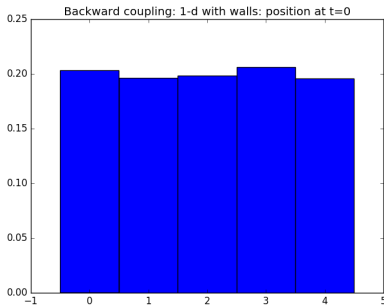
- Dictionary of random maps going back in time.

Coupling from the past 6/8



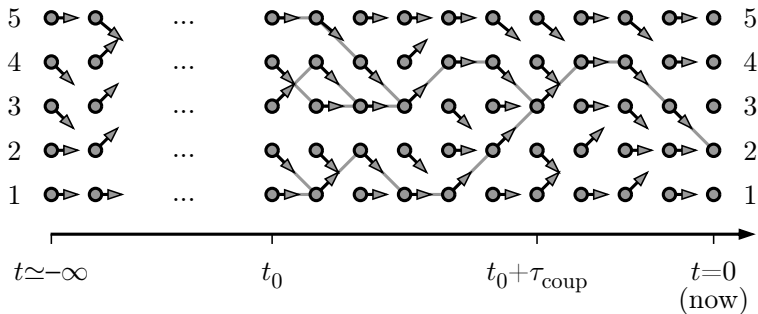
- Starting an MCMC simulation at $t = -\infty$
- Propp & Wilson (1997)

Coupling from the past 7/8



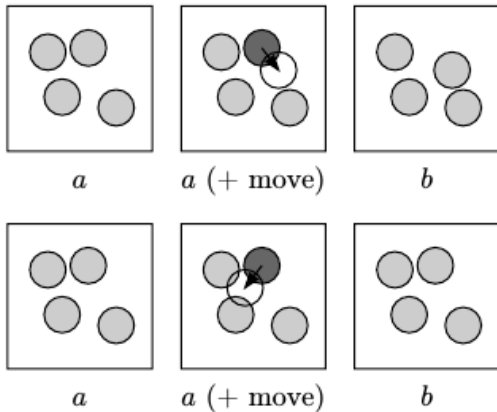
- Perfect sample at $t = 0$, starting from $t = -\infty$
- Propp & Wilson (1997)

Coupling from the past 8/8

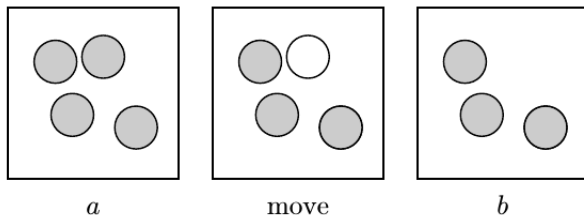


- Try it yourself!

Hard-sphere simulation (traditional)



Hard-sphere simulation (birth-and-death)



$$Z = \sum_{N=0}^{\infty} \lambda^N \int \cdots \int dx_1 \dots dx_N \pi(x_1, \dots, x_N)$$

- $\pi(a) = \lambda \pi(b)$
- Death probability (per particle, per time interval): $1dt$
- Birth probability (per unit square): λdt

Poisson distribution

Poisson distribution (number n of events per unit time):

$$\pi_{\Delta t=1}(n) = \frac{\lambda^n e^{-\lambda}}{n!}$$

Poisson distribution (number n of events per time dt):

$$\pi_{dt}(n) = \frac{(\lambda dt)^n e^{-\lambda dt}}{n!} \implies \pi_{dt}(1) = \lambda dt, \pi_{dt}(2) = 0$$

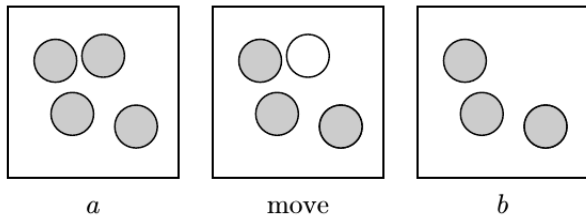
Poisson waiting time: Probability that next event after time t :

$$\mathbb{P}(t) = (1 - \lambda dt), \dots, (1 - \lambda dt) \lambda dt$$

$$\mathbb{P}(t) = \underbrace{\left(\overbrace{(1 - \lambda dt) \rightarrow (1 - \lambda dt)}^{\sum dt=t} \right)}_{e^{-\lambda t}} \lambda dt$$

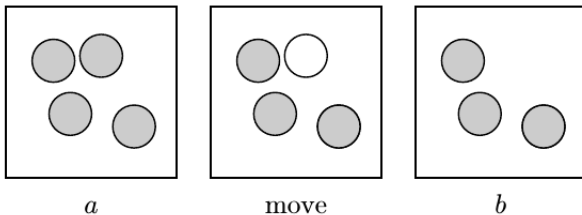
...can be sampled with $t = (-\log \text{ran}[0, 1])/\lambda$

Birth-and-death (principle 1)



- N spheres, each of them may die.
- a new sphere may be born (but there may be problems).
- rate for next event: $N + \lambda$.
- $\mathbb{P}(\text{death}) \propto N$ and $\mathbb{P}(\text{birth}) \propto \lambda$, reject if overlap.

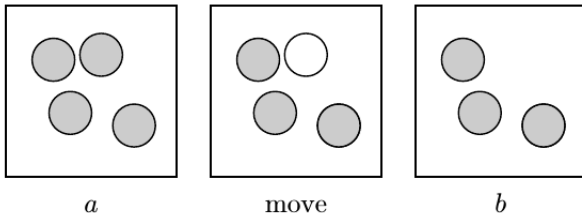
Birth-and-death (implementation 1)



- start with $N = 0$ spheres
- Go to next-event time : $-\log \text{ran}/(N + \lambda)$ (in steps of 1)
- sample random number $\text{ran}[0, 1]$: if smaller than $\lambda/(\lambda + N)$:
add a disk (reject if overlap), otherwise delete a disk.

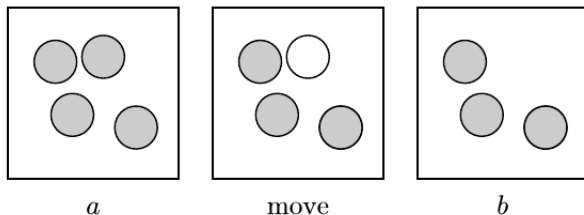
NB: Check configuration at integer time steps, for sampling.

Birth-and-death (principle 2)



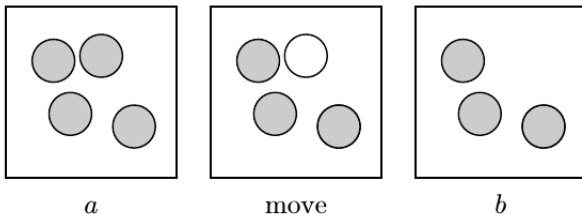
- N spheres, each of them knows when it will die (sad) rate=1.
- a new sphere may be born (but there may be problems) rate $= \lambda$.

Birth-and-death (implementation 2)



- start with $N = 0$ spheres.
- Advance to next birth time : $-\log \text{ran}[0, 1]/\lambda$ (in steps of 1).
- If no rejection, install death time $-\log \text{ran}[0, 1]$

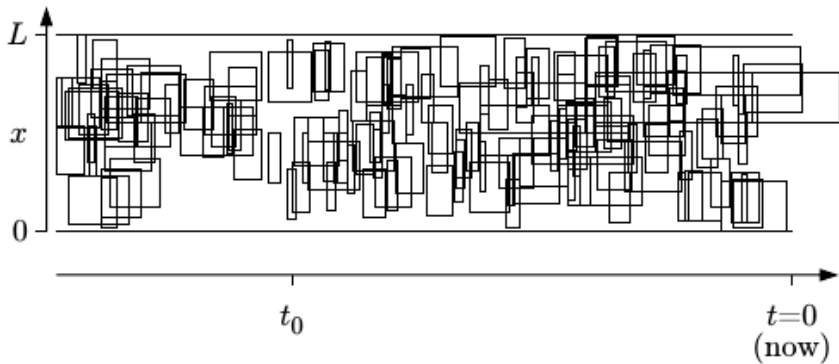
Birth-and-death (principle 3)



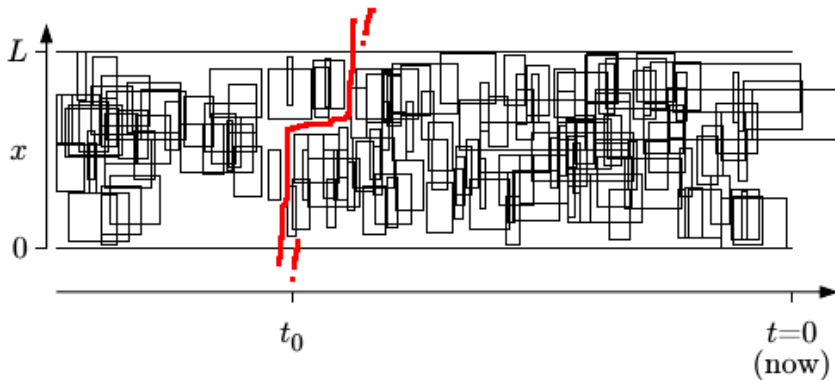
- Hypothetical spheres are born with rate $= \lambda$, and they die with rate 1.

Check later whether all this pans out correctly.

Birth-and-death (implementation 3)

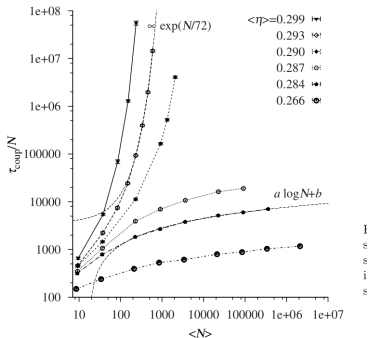


Birth-and-death (implementation 3)



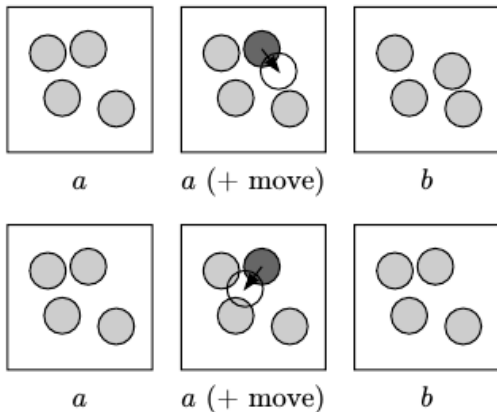
- Can be made into a perfect sampling algorithm
- Wilson (2000)

Birth-and-death (implementation 3)



- Bernard et al. (2010)
- Dynamical phase transition

Hard-sphere simulation (traditional)



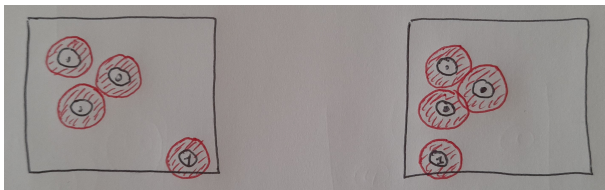
Algorithm remains correct if displacement random in box.



- At low density, any two configurations of spheres a and z can be connected through a path of length $< 2N$ as follows:
 $a \rightarrow b \rightarrow c \rightarrow \dots \rightarrow z$, where any two neighbors differ only in 1 sphere.
- MC algorithm: Take random sphere, place it at random position anywhere in the box.

Kannan et al. (2003)

Path coupling 2/4



- MC algorithm: Take random sphere, place it at the same random position for both copies.
- $p(1 \rightarrow 0)$: Pick 1, move to where it fits in both copies

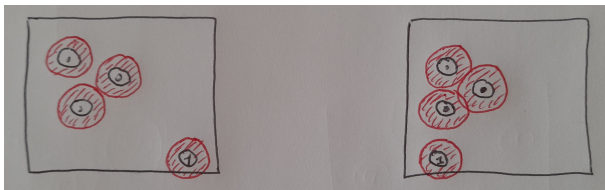
$$p(1 \rightarrow 0) \geq \frac{1}{N} \left[1 - \frac{N-1}{N} 4\eta \right]$$

- $p(1 \rightarrow 2)$: Pick $2 \dots N$ move near to 1_A or 1_B .

$$p(1 \rightarrow 2) \leq \frac{N-1}{N} \left[\frac{8}{N} \eta \right]$$

- \implies for $\eta < 1/12$: further coupling likely.

Path coupling 3/4



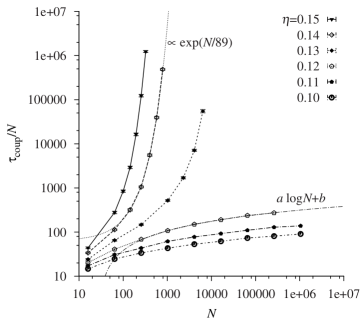
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- \implies for $\eta < 1/12$: further coupling likely.



- Bernard et al. (2010)
- Damage-spreading dynamical phase transition

Helmuth et al. (2020)

Strategies for overcoming the limitations of MCMC

- Larger moves—faster convergence
- Exact-sampling approaches from MCMC