# Algorithms and computations in physics (Oxford Lectures 2025)

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Version 20 Jan 2025

We start our parallel exploration of physics and of computing with the concept of sampling, the process of producing examples ("samples") of a probability distribution. In week 1, we consider "direct" sampling (the examples are obtained directly) and, among the many connections to physics, will come across the Maxwell distribution. In 1859, it marked the beginning of the field of statistical physics.

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# 1 Direct sampling

In this first lecture, I introduce to the concept of sampling. Some of the material can be found in my textbook [2]. For background on statistics, one may consult the book by Wasserman [1], among others.

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Werner Krauth: Algorithms and Computations in Physics (2025 Oxford Lectures)

#### 1.1 Children on the beach in Monaco

Direct sampling <sup>1</sup> is exemplified by a children's game popular in Monaco, south of France, whence the name "Monte Carlo" method. Children fling pebbles (small stones) into a square. They then count how many of them fall into the circle inscribed in the square (see [2, Sect. 1.1] for the full story).

#### 1.1.1 Sample spaces and random variables



Figure 1: Children on the Monte Carlo beach, computing the number  $\pi$ .

The square in Fig. 1 constitutes a sample space  $\Omega$ . <sup>2</sup> Modern probability theory, since Kolmogoroff, then assigns probabilities, not in general to individual samples (x, y), but to subsets of the sample space  $\Omega$ . A small subset of size dxdy around a point (x, y) then has the probability

$$\pi^{\Box}(x,y)\mathrm{d}x\mathrm{d}y = \begin{cases} \mathrm{const} \, \mathrm{d}x\mathrm{d}y & \text{if } x, y \in \Omega^{\Box} \\ 0 & \text{else} \end{cases}.$$
 (1)

where, in this case, we know that const = 1/4. The pebbles that the children throw are (exact, perfect) samples of the distribution  $\pi$ , and the act of throwing the pebbles is denoted as "sampling (from) the distribution  $\pi$ ".

On the Monte Carlo beach, a circle is inscribed inside the square. While all pebbles inside the square are referred to as *trials*, those inside the circle are *hits*. The function on the sample space defines a "random variable" <sup>3</sup>.  $\mathcal{O} : \Omega \to \mathbb{R}$ :

$$\mathcal{O}[(x,y)] = \begin{cases} 1 & \text{if } (x,y) \in \odot \\ 0 & \text{else} \end{cases}$$
(2)

The quantity  $\mathcal{O}$  is a "Bernoulli random variable" that we can also sample, more directly, as

$$\mathcal{O}[(x,y)] = \begin{cases} 1 & P = \pi/4 \\ 0 & P = 1 - \pi/4 \end{cases},$$
(3)

but then we need to know the value of  $\pi$ . The passage from eq. (2) to eq. (3) illustrates that sample spaces tend to disappear in discussions of probability theory [1, p. 27]. If we have forgotten the value of  $\pi$ , we may estimate it from the pebble throws. This moves us from probability theory into the field of statistics.

<sup>&</sup>lt;sup>1</sup>The word "sample" has the same origin as "example"

<sup>&</sup>lt;sup>2</sup>Sample space: the set of all outcomes of the pebble-throw experiments

<sup>&</sup>lt;sup>3</sup>random comes from an old French word that still survives in randonnée, in other words a hike

#### 1.1.2 Pseudocode, pseudo-random numbers

The children's game is the first of 150 algorithms that we will discuss in the present lecture series. We will specify them through self-explanatory pseudo-code (see Alg. 1 (direct-pi)). In Alg. 1 (direct-pi), as throughout this course, we take for granted the existence of uniform

```
 \begin{array}{l} \textbf{procedure direct-pi} \\ N_{\rm hits} \leftarrow 0 \ (\rm initialize) \\ \textbf{for } i = 1, \ldots, N \textbf{:} \\ \left\{ \begin{array}{l} x \leftarrow {\tt ran}(-1,1) \\ y \leftarrow {\tt ran}(-1,1) \\ \textbf{if } x^2 + y^2 < 1 \ N_{\rm hits} \leftarrow N_{\rm hits} + 1 \\ \textbf{output } N_{\rm hits} \end{array} \right. \end{array}
```

Algorithm 1: direct-pi. Using the children's game with N pebbles to compute  $\pi$ .

random numbers ran(a, b) in the interval between a and b. During 80 years, now, such numbers have been generated as pseudo-random numbers, in other words deterministic procedures that look random, and that condense deep concepts from number theory and cryptography. Up to the year 2000, roughly, problems with random numbers were common, and users on high alert, running batteries of checks on their results for different classes of random numbers. In modern times, the quality of random-number generators has much improved, although they continue to have flaws and will themselves be superseded by even more sophisticated variants. What will remain is the pseudo-code, as in Algorithm 1, and we will stick to it.

#### 1.1.3 From Punch-card coding to Python and to ChatGPT

The pseudocode of Alg. 1 (direct-pi), and of all algorithms that follow must be communicated to a computer. This, over the decades, first consisted in punching holes into cards, of which one had huge piles. Later, pseudocode was to be translated into one or the other computer language, with Python becoming more and more popular, although the language itself changes over time. Today, pseudo-code as Alg. 1 (direct-pi) may be simply copied and pasted into a window of an AI program See this ChatGPT output. This is extremely helpful, and it may make interfacing with a computer appear as effortless as flinging  $4 \times 10^8$  pebbles on a modern laptop rather than on the Monte Carlo beach. Nevertheless, the concept of effortless creation certainly remains a mirage. In whichever way one goes about it, actual running the short example programs of this course will be a powerful way to confront oneself with its subjects. As this has become so easy, why not implement them all?

Run	$N_{\rm hits}$	Estimate of $\pi$
1	3156	3.156
2	3150	3.150
3	3127	3.127
4	3171	3.171
5	3148	3.148

**Table 1:** Results of five runs of Alg. 1 (direct-pi) with N = 4000. Approximations of  $\pi$  are obtained by shifting a decimal point.

Implementing Alg. 1 (direct-pi), not in a computer, but on the beach, the children record the proportion of "hits", that is, of the fraction of pebbles inside the circle (see Table 1.1.3, for N = 4000). Without knowing it, they have computed a two-dimensional integral:

$$\underbrace{\frac{N_{\text{hits}}}{\text{trials}} = \frac{1}{N} \sum_{i=1}^{N} \mathcal{O}_i}_{\text{sampling}} \simeq \underbrace{\langle \mathcal{O} \rangle = \frac{\int_{-1}^1 dx \int_{-1}^1 dy \ \pi(x,y) \mathcal{O}(x,y)}{\int_{-1}^1 dx \int_{-1}^1 dy \ \pi(x,y)}}_{\text{integration}}.$$
(4)

On the left of eq. (4), there is the "running average" of the random variable  $\mathcal{O}$ , evaluated at the samples  $(x_i, y_i)$ . The probability distribution  $\pi(x, y)$  is absent, rather than being evaluated, it is sampled. By virtue of the law of large numbers, this running average converges (in a way we will need to discuss precisely, later) to the expectation (mean value) of  $\mathcal{O}$ , expressed as a ratio of two two-dimensional integrals. The same approach allows one to evaluate (that is, to sample) high-dimensional integrals in physics and other domains, if only we can think of how to do the sampling. What would happen in the limit  $N \to \infty$  is easy to grasp, but we will later pursue the more relevant question what precisely we learn from 3156 hits for 4000 trials.

#### 1.2 Direct sampling—discrete distributions

Sampling defines the Monte Carlo method (and some other fields, like machine learning), and direct sampling is the first step on this journey. In this section, we treat the case of a discrete distribution in a sample space  $\Omega = \{1, \ldots, K\}$ , with non-uniform weights  $\{\pi_1, \ldots, \pi_K\}$ . The sampling is non-trivial already for moderate values of K, as we experience ourselves each week in the "Saturday night problem", when the sample 1 and its probability  $\pi_1$  correspond to studying (this course), 2 corresponds to cleaning the house, 3 to getting some exercise, and so on. There are only a few choices but, clearly, it takes us so long to decide what to do, after all.

The optimal way to go about sampling a finite distribution has not been known for a very long time [3], but we will first treat the standard approaches, rejection sampling, and tower sampling, which are interesting in their own right.

#### 1.2.1 Rejection sampling

Rejection sampling is the simplest approach one can think of: place all the items into a big box such that they do not overlap, then throw pebbles into the box and stop when one of them hits an item. It helps to make the box rectangular, and to place the items side by side (see Alg. 2 (reject-finite)).

```
procedure reject-finite

\pi_{\max} \leftarrow \max_{k=1}^{K} \pi_k

1 k \leftarrow \operatorname{nran}(1, K)

\Upsilon \leftarrow \operatorname{ran}(0, \pi_{\max})

if \Upsilon > \pi_k goto 1

output k
```

Algorithm 2: reject-finite. Sampling a finite distribution  $\{\pi_1, \ldots, \pi_K\}$  with a rejection algorithm.



Figure 2: Saturday night problem solved by Alg. 2 (reject-finite).

#### 1.2.2 Tower sampling

In what we call "tower sampling", the K probabilities are stacked onto each other, and all the probabilities are added up as  $\Pi_0 = 0$ ,  $\Pi_1 = \pi_1$ ,  $\Pi_2 = \pi_1 + \pi_2$ , and so on. Then we throw a uniform pebble  $ran(0, \Pi_K)$ .

```
procedure tower-sample

input \{\pi_1, \dots, \pi_K\}

\Pi_0 \leftarrow 0

for l = 1, \dots, K: \Pi_l \leftarrow \Pi_{l-1} + \pi_l

\Upsilon \leftarrow \operatorname{ran}(0, \Pi_K)

* find k with \Pi_{k-1} < \Upsilon < \Pi_k

output k
```

Algorithm 3: tower-sample. Tower sampling of a finite distribution  $\{\pi_1, \ldots, \pi_K\}$  without rejections. The search indicated in the line marked with a \* makes its complexity to be  $\mathcal{O}(\log K)$  per sample, once we have constructed the tower, in  $\mathcal{O}(K)$  steps.

Tower sampling can be applied to discrete distributions with a total number K in the hundreds, thousands, or even millions. It is used when the naive rejection method of Fig. 2 fails because of too many rejections. Tower sampling becomes impracticable only when the probabilities  $\{\pi_1, \ldots, \pi_K\}$  can no longer be listed. This rejection-free method is not optimal, but easy to implement and of theoretical interest.



Figure 3: Saturday night problem solved by tower sampling.

#### 1.2.3 Complexity of discrete sampling—Walker's method

Suppose we have  $\pi_1, \ldots, \pi_N$ , with N much larger than 1. What is the complexity of sampling the distribution  $\{\pi_1, \ldots, p_N\}$ ? The tower method of Sec. ?? is of complexity  $\mathcal{O}(\log N)$ , but this is not optimal. An extremely useful algorithm called Walker's method of aliases [3], proves that the complexity of sampling an  $x \sim \pi$  is  $\mathcal{O}(1)$ , that is, it is independent of N. In the Saturday-night problems, if you had a billion 10<sup>9</sup> activities to choose from, it would take us only  $\mathcal{O}(1)$  operations to find out which one to choose. (There is an initialization step of complexity  $\mathcal{O}(10^9)$ , which limits the applicability of the method.) As illustrated in Fig. 4, we represent the probabilities  $\pi_i$  as "slabs", then put these slabs into two buckets, one containing the small ones and the other the large ones. Then we take a small one, and put a large one on top, then cut off at the mean value....



**Figure 4:** Walker's algorithm. (a): Distribution  $\{\pi_1, \ldots, \pi_N\}$  to be sampled. Each  $\pi_i$  is represented by a slab. The mean of the  $\pi_i$  is shown. (b): Rearrangement of the slabs into a perfect rectangle, from which the sampling is straightforward.

In a few weeks, we will implement discrete sampling for the Ising model, say, on a  $1000 \times 1000$  lattice in two dimensions, where  $|\Omega| = 2^{1000000} = 9.9 \times 10^{301029}$ , where Walker's method cannot be used. But this is for a special choice of  $\{\pi_1, \ldots, \pi_N\}$ .

Walker's method, as discussed, samples *one* element, k, of  $\{1, \ldots, N\}$ , with probability  $\pi_k$ . A related problem consists in sampling a subset of  $\{1, \ldots, N\}$  (an element of the powerset of  $\{1, \ldots, N\}$ , where we chose each element k with probability  $\pi_k$ . This can also be solved with Walker's method, but in a distorted way.

## 1.3 Direct sampling—continuous distributions

We now consider the sampling problem for continuous distributions, but start by approaching the continuum limit from the aforementioned discretized problem. We discover a deep relation between integration and sampling, and between the substitution of variables and what we call a *sample transformation*.

#### **1.3.1** Sample transformation—simple examples



Figure 5: Tower sampling for a discretized version of  $\pi(x) = (\gamma + 1)x^{\gamma}$  in the interval  $x \in (0, 1]$ . The case  $\gamma = -\frac{1}{2}$  is shown).

We consider the continuum limit of tower sampling. As an example, let us sample random numbers 0 < x < 1 distributed according to an algebraic function  $\pi(x) = (\gamma + 1)x^{\gamma}$  (with  $\gamma > -1$ ) (see Fig. 5). A pebble in the tower must be identified with its corresponding x position. In the continuum limit, we find

$$\pi(x) = (\gamma + 1)x^{\gamma} \text{ for } 0 < x < 1,$$
  

$$\Pi(x) = \int_0^x dx \ \pi(x') = x^{\gamma+1} = \operatorname{ran}(0, 1),$$
  

$$x = \operatorname{ran}(0, 1)^{1/(\gamma+1)}.$$
(5)

The transformation method corresponds to a sample transformation: For the above algebraic function, we can transform the integral over a flat distribution into the integral of the target distribution:

$$\int_0^1 \mathrm{d}\Upsilon \quad \xrightarrow{\text{integral}}_{\text{transform}} \text{ const} \int_0^1 \mathrm{d}x \ x^\gamma.$$

We now treat  $\Upsilon$  ("Upsilon") as a sample  $\Upsilon = ran(0, 1)$ , and it transforms as follows:

 $\mathrm{d}\Upsilon = \mathrm{const} \cdot \mathrm{d}x \ x^{\gamma}.$ 

$$\operatorname{ran}(0,1) = \Upsilon = \operatorname{const}' \cdot x^{\gamma+1} + \operatorname{const}''.$$

Finally (checking that the bounds of ran(0, 1) correspond to x = 0 and x = 1), this results in

$$x = \operatorname{ran}(0, 1)^{1/(\gamma+1)}, \tag{6}$$

in agreement with eq. (5).

As a second example of sample transformation, we consider exponentially distributed random numbers, so that  $\pi(x) \propto e^{-\lambda x}$  for  $x \ge 0$ . We again write

$$\int_0^1 \mathrm{d}\Upsilon = \mathrm{const} \ \int_0^\infty \mathrm{d}x \ \mathrm{e}^{-\lambda x} \tag{7}$$

and <u>see</u> on the left-hand side of eq. (7), that  $\Upsilon = ran(0, 1)$ , so that

$$\begin{split} \mathrm{d} \Upsilon &= \mathrm{const} \cdot \mathrm{d} x \ \mathrm{e}^{-\lambda x},\\ \mathrm{ran}(0,1) &= \Upsilon = \mathrm{const}' \cdot \mathrm{e}^{-\lambda x} + \mathrm{const}''. \end{split}$$

Checking the bounds x = 0 and  $x = \infty$ , this leads to

$$-\frac{1}{\lambda}\log\operatorname{ran}(0,1) = x. \tag{8}$$

Each "negative logarithm of ran(0,1)" will from now be recognized as an exponential random number.

#### 1.3.2 Sample transformation—Gaussian, multidimensional Gaussians

In our exploration of sample transformations, we move ahead to more complex, and physically relevant, cases involving Gaussian random numbers x that are a staple in many fields of science. Restricting ourselves to unit variance  $\sigma^2 = 1$ , they are distributed as

$$\pi(x) = \frac{1}{\sqrt{2\pi}} \exp\left[-\frac{x^2}{2}\right].$$

Subroutines for Gaussian random numbers are readily available, but let us look under the hood of the corresponding algorithms. Gaussians again illustrate sample transformation.

To evaluate the error integral

$$\int_{-\infty}^{\infty} \frac{\mathrm{d}x}{\sqrt{2\pi}} \mathrm{e}^{-x^2/2} = 1,\tag{9}$$

we recall that we should square eq. (9):

$$\left[\int_{-\infty}^{\infty} \frac{\mathrm{d}x}{\sqrt{2\pi}} \exp\left(-x^2/2\right)\right]^2 = \int_{-\infty}^{\infty} \frac{\mathrm{d}x}{\sqrt{2\pi}} \mathrm{e}^{-x^2/2} \int_{-\infty}^{\infty} \frac{\mathrm{d}y}{\sqrt{2\pi}} \mathrm{e}^{-y^2/2} \tag{10}$$

$$= \int_{-\infty}^{\infty} \frac{\mathrm{d}x \, \mathrm{d}y}{2\pi} \exp\left[-(x^2 + y^2)/2\right],\tag{11}$$

introduce polar coordinates  $(dx dy = rdr d\phi)$ ,

$$\ldots = \int_0^{2\pi} \frac{\mathrm{d}\phi}{2\pi} \int_0^\infty r \,\mathrm{d}r \,\exp\left(-r^2/2\right),$$

and finally substitute  $r^2/2 = \Upsilon ~(r \, \mathrm{d} r ~= \mathrm{d} \Upsilon$  )

$$\dots = \underbrace{\int_{0}^{2\pi} \frac{\mathrm{d}\phi}{2\pi}}_{1} \underbrace{\int_{0}^{\infty} \mathrm{d}\Upsilon \ \mathrm{e}^{-\Upsilon}}_{1} \cdot \underbrace{(12)}_{1}$$

Equation (12) famously implies eq. (9), but it also shows us the way to obtaining independent Gaussian samples x and y. It suffices to sample the uniform distribution of  $\phi$  and the exponential distribution for  $\Upsilon$ , as indicated, and then to transform everything back to have two independent Gaussian random variables x and y (see Alg. 4 (gauss), the computation of sines and cosines can be avoided [2, Sect. 1.2.5]).

.

procedure gauss  
input 
$$\sigma$$
  
 $\phi \leftarrow \operatorname{ran}(0, 2\pi)$   
 $\Upsilon \leftarrow -\log \operatorname{ran}(0, 1)$   
 $r \leftarrow \sigma \sqrt{2\Upsilon}$   
 $x \leftarrow r \cos \phi$   
 $y \leftarrow r \sin \phi$   
output  $\{x, y\}$ 

Algorithm 4: gauss. Two independent Gaussian random numbers obtained by sample transformation.

As we just learned to sample the one-dimensional Gaussian integral, we can also sample the *d*-dimensional Gaussian integral,

$$1 = \int \dots \int \underbrace{\mathrm{d}x_1 \dots \mathrm{d}x_d}_{\mathrm{d}V} \left(\frac{1}{\sqrt{2\pi}}\right)^d \exp\left[-\frac{1}{2}(x_1^2 + \dots + x_d^2)\right],\tag{13}$$

by running Alg. 4 (gauss) d/2 times. This obtains  $\{x_1, \ldots, x_d\}$ . We now substitute

$$\mathrm{d}x_1 \dots \mathrm{d}x_d = r^{d-1} \,\mathrm{d}r \,\mathrm{d}\Omega$$

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to obtain:

$$1 = \left(\frac{1}{\sqrt{2\pi}}\right)^d \underbrace{\int_0^\infty \mathrm{d}r \ r^{d-1} \exp\left(-r^2/2\right)}_{\text{uniform solid angle}} \underbrace{\int_0^\infty \mathrm{d}\Omega}_{\text{uniform solid angle}}.$$
 (14)

The  $\{x_1, \ldots, x_N\}$  are samples, and so is  $r = \sqrt{x_1^2 + \cdots + x_N^2}$ , and so is the solid angle  $\Omega^4$ . With a little rearrangement of the radius, we either obtain uniformly distributed pebbles within the unit *d*-dimensional hypersphere (Algorithm 5) or on its surface (Algorithm 6, see also Fig. 6).

```
procedure direct-sphere

\Sigma \leftarrow 0

for k = 1, ..., d:

\begin{cases} x_k \leftarrow gauss(\sigma) \\ \Sigma \leftarrow \Sigma + x_k^2 \end{cases}

\Upsilon \leftarrow ran(0, 1)^{1/d}

for k = 1, ..., d:

\{ x_k \leftarrow \Upsilon x_k / \sqrt{\Sigma} \end{cases}

output \{x_1, ..., x_d\}
```

Algorithm 5: direct-sphere. Uniform random vector inside the *d*-dimensional unit sphere. The output is independent of  $\sigma$ .

```
\begin{array}{l} \textbf{procedure direct-surface}\\ \sigma\leftarrow 1/\sqrt{d}\\ \Sigma\leftarrow 0\\ \textbf{for }k=1,\ldots,d\textbf{:}\\ \left\{\begin{array}{l} x_k\leftarrow \textbf{gauss}(\sigma)\\ \Sigma\leftarrow \Sigma+x_k^2\\ \textbf{for }k=1,\ldots,d\textbf{:}\\ \left\{\begin{array}{l} x_k\leftarrow x_k/\sqrt{\Sigma}\\ \textbf{output }\{x_1,\ldots,x_d\}\end{array}\right.\end{array}
```

Algorithm 6: direct-surface. Random vector on the surface of the *d*-dimensional unit sphere. For large d,  $\Sigma$  approaches one (see Fig. 6).



Figure 6: Random samples on the surface of the 3-dimensional sphere, from Alg. 6 (direct-surface)

<sup>&</sup>lt;sup>4</sup>not to be confused with the sample space

The Gaussian is unique in combining into a *d*-dimensional isotropic integral. The transformation from eq. (13) to eq. (14), that is, Alg. 6 (direct-surface), was already known to Maxwell, and was the key to his invention of the Maxwell distribution. Later generalized by Boltzmann, it became the core of all of statistical mechanics. So, maybe it was worth looking under the hood of the Gaussian algorithm.

#### 1.3.3 Rejection method

We have a closer look at rejection sampling, after Sec. 1.2.1, but now for a continuum distribution. In preparation of Lecture 2 and following Ref. [4], we consider the Boltzmann distribution of the anharmonic oscillator

$$\int_{-\infty}^{\infty} \mathrm{d}x \underbrace{\exp\left(-\frac{x^2}{2} - \frac{x^4}{4}\right)}_{\pi_{24}(x)},\tag{15}$$

that we need not normalize, and likewise for the Gaussian  $\pi_2$ .



Figure 7: Naive and state-of-the-art algorithms to sample the Boltzmann distribution for the anharmonic oscillator as uniformly distributed pebbles. (a): The naive algorithm introduces a cutoff  $x_{\max}$  and discards many pebbles. (b): Sampling pebbles (x, y) below the Gaussian curve and then discarding those samples (x, y) above  $\pi_{24}(x)$  is a winning strategy (see eq. (17)

It is straightforward to enclose  $\pi_{24}$  in a rectangular box between  $x = \pm x_{\text{max}}$  and to adapt Alg. 2 (reject-finite) (see Fig. 7a). But this is wasteful of pebbles and potentially dangerous because of the cutoff. The box in Fig. 7a is nothing but a bounding function that satisfies two conditions: it can be sampled and it dominates the distribution of interest. For the anharmonic oscillator of eq. (15), the Gaussian is an alternative bounding function, so that we can write:

$$\int_{-\infty}^{\infty} dx \, \exp\left(-\frac{x^2}{2} - \frac{x^4}{4}\right) = \underbrace{\int_{-\infty}^{\infty} dx \, \exp\left(-\frac{x^2}{2}\right)}_{\text{sample}} \underbrace{\exp\left(-\frac{x^4}{4}\right)}_{\text{accept/reject}} \tag{16}$$

This suggests that to sample the anharmonic oscillator, we may sample a Gaussian, and then accept the pebble x with probability  $\exp\left(-x^4/4\right)$  (see Alg. 7 (direct-anharm)). But to convince us that this brashly introduced algorithm is actually correct, we sample the Gaussian in x, then spread pebbles out evenly on the y-axis with a  $\operatorname{ran}(0, \exp\left(-x^2/2\right))$ . We now have uniform pebbles in the two-dimensional region below the Gaussian. It suffices to reject all pebbles above  $\pi_{24}$ , and this is what Alg. 7 (direct-anharm) implements.

To <u>see</u> how to proceed in general, we write:

$$\underbrace{\int_{-\infty}^{\infty} dx \, \exp\left(-\frac{x^2}{2}\right)}_{\text{sample}} \underbrace{\exp\left(-\frac{x^4}{4}\right)}_{\text{accept/reject}} = \underbrace{\int_{-\infty}^{\infty} dx \, \exp\left(-\frac{x^2}{2}\right)}_{\text{sample } x} \underbrace{\frac{\exp\left(-\frac{x^2}{2} - \frac{x^4}{4}\right)}_{\text{sample } y, \text{ see Fig. 7b}}.$$
 (17)

In conclusion, to sample a distribution, such as  $\pi_{24}$ , we may divide by and multiply it with another distribution, such as  $\pi_2$ , under the two conditions that we know to sample the latter, and that it is always larger than the former. Remarkably, neither distribution (neither  $\pi_2$  nor  $\pi_{24}$ ) need be normalized. The procedure works if the proportion of pebbles we must discard is not too large.

```
procedure direct-anharm
while True:
\begin{cases} x \leftarrow gauss(0, 1/\sqrt{\beta}) \\ y \leftarrow ran[0, \pi_2(x)] \\ if \ y < \pi_{24}(x) \text{ break} \\ output \ x \end{cases}
```

Algorithm 7: direct-anharm. Sampling  $\pi_{24}$  through the rejection of Gaussians samples from Eq. (17).

#### 1.4 Direct sampling: fundamental aspects

Many fundamental aspects of sampling already manifest themselves in the direct-sampling framework, and then translate, *mutatis mutandis*, to the much more complicated Markov chains of subsequent lectures. For example, the strong law of large numbers that we need to understand for direct samples will turn into the famous ergodic theorem for Markov chains. We also discuss importance sampling that permeates all of Monte Carlo and discuss the frequentist interpretation of probabilities at the core of the method.

#### 1.4.1 Laws of large numbers—importance sampling

To discuss the convergence of Markov chains, we consider a frankly difficult sampling problem, the  $\gamma$  integral:

$$I(\gamma) = \int_0^1 \mathrm{d}x \ x^\gamma = \frac{1}{\gamma+1} \quad \text{for } \gamma > -1 \tag{18}$$

(see [2, Sect. 1.4.2] for the full context). We attempt to compute the integral in a sample space  $\Omega^{[0,1]}$ , the unit interval between 0 and 1.

$$I(\gamma) = \int_0^1 \mathrm{d}x \ x^\gamma = \int_0^1 \underbrace{(\mathrm{1d}x)}_{x=\mathrm{ran}(0,1)} \underbrace{\overset{\mathcal{O}}{x^\gamma}}_{(19)}$$

As we discussed before, the random variable  $\mathcal{O}$  has its own probability distribution:

$$\pi(\mathcal{O}) = (\alpha - 1)\mathcal{O}^{-\alpha},\tag{20}$$

with  $\alpha = 1 - 1/\gamma$ . Its mean value of the random variable  $\mathcal{O}$  can be equivalently written with  $\pi(\mathcal{O})$  or in the original sample space:

$$\langle \mathcal{O} \rangle = (\alpha - 1) \int_{1}^{\infty} \mathrm{d}\mathcal{O} \ \mathcal{O}\mathcal{O}^{-\alpha} = \int_{0}^{1} \mathrm{d}x x^{\gamma}.$$
 (21)

The same holds for any higher moments.

After these preliminaries, let us now actually compute the  $\gamma$  integral with a running average of a sum of uniform random numbers to the power of  $\gamma$  (see Algorithm 8). We compute the integral  $I(\gamma)$  as the mean value of  $x^{\gamma}$ , with  $x = \operatorname{ran}(0, 1)$ , and likewise "try" to compute the error through the Gaussian error formula

$$\operatorname{Error} \stackrel{?}{=} \frac{\sqrt{\langle \mathcal{O}^2 \rangle - \langle \mathcal{O} \rangle^2}}{\sqrt{N}} \tag{22}$$

This is implemented in Alg. 8 (direct-gamma) (which we should modify to output the Gaussian error). The calculation works well for large enough  $\gamma \gtrsim -0.5$  but seems to fail for  $-1 < \gamma < -0.5$ .

 $\begin{array}{l} \textbf{procedure direct-gamma} \\ \Sigma \leftarrow 0 \\ \textbf{for } i = 1, \ldots, N \textbf{:} \\ \left\{ \begin{array}{l} x_i \leftarrow \texttt{ran}(0, 1) \\ \Sigma \leftarrow \Sigma + x_i^{\gamma} \text{ (running average: } \Sigma/i) \\ \textbf{output } \Sigma/N \end{array} \right. \end{array}$ 

Algorithm 8: direct-gamma. Computing the  $\gamma$ -integral in eq. (18) by direct sampling.

$\gamma$	$\Sigma/N \pm \text{Error}$	$1/(\gamma + 1)$
2.0	$0.334 \pm 0.003$	$0.333\ldots$
1.0	$0.501 \pm 0.003$	0.5
0.0	$1.000\pm0.000$	1
-0.2	$1.249\pm0.003$	1.25
-0.4	$1.682\pm0.014$	1.666
-0.8	$3.959 \pm 0.110$	5.0

**Table 2:** Output of Alg. 8 (direct-gamma) for various values of  $\gamma$  ( $N = 10\,000$ , standard empirical error shown). The computation for  $\gamma = -0.8$  is in trouble.

$\gamma$	$\zeta$	$\Sigma/N$	$\frac{\zeta+1}{\gamma+1}$
-0.4	0.0	$1.685\pm0.017$	1.66
-0.6	-0.4	$1.495\pm0.008$	1.5
-0.7	-0.6	$1.331\pm0.004$	1.33
-0.8	-0.7	$1.508\pm0.008$	1.5

**Table 3:** Output of Alg. 9 (direct-gamma-zeta) with  $N = 10\,000$ . All pairs  $\{\gamma, \zeta\}$  satisfy  $2\gamma - \zeta > -1$  so that  $\langle \mathcal{O}^2 \rangle < \infty$ .

The problem has to do that for  $\gamma < -\frac{1}{2}$ , the variance of  $\mathcal{O}$  is infinite, so that the Gaussian error analysis no longer applies. A modified program, Alg. 9 (direct-gamma-zeta), uses importance



Figure 8: Running average of Alg. 8 (direct-gamma) for  $\gamma = -0.8$ . The strong law of large numbers guarantees that  $\Sigma_i/i$  converges towards 5, that is, that for any  $\epsilon$ , there is an  $i_{\epsilon}$  such that  $|\Sigma_j/j - 5| < \epsilon \forall j > i_{\epsilon}$ .

sampling to reduce the non-Gaussian fluctuations. It computes, not  $I(\gamma)$ , but  $I(\gamma)/I(\zeta)$ :

$$\Sigma/N = \frac{1}{N} \sum_{i=1}^{N} \mathcal{O}_i \simeq \langle \mathcal{O} \rangle = \frac{\int_0^1 \mathrm{d}x \ \pi(x)\mathcal{O}(x)}{\int_0^1 \mathrm{d}x \ \pi(x)} = \frac{\int_0^1 \mathrm{d}x \ x^{\zeta} x^{\gamma-\zeta}}{\int_0^1 \mathrm{d}x \ x^{\zeta}} = \frac{\int_0^1 \mathrm{d}x \ x^{\gamma}}{\int_0^1 \mathrm{d}x \ x^{\zeta}} = \frac{I(\gamma)}{I(\zeta)} = \frac{\zeta+1}{\gamma+1}.$$
 (23)

(see [2, eq. 1.74] for details).

procedure direct-gamma-zeta  $\Sigma \leftarrow 0$ for i = 1, ..., N:  $\begin{cases} x_i \leftarrow \operatorname{ran}(0, 1)^{1/(\zeta+1)} & (\pi(x_i) \propto x_i^{\zeta}, \text{ see eq. (6)}) \\ \Sigma \leftarrow \Sigma + x_i^{\gamma-\zeta} \\ \text{output } \Sigma/N \end{cases}$ 

Algorithm 9: direct-gamma-zeta. Using importance sampling to compute the  $\gamma$ -integral (see eq. (23)).

## 1.4.2 What we learn about $\pi$ by throwing pebbles (not treated in week 1)

What can we learn (about  $\pi$ ) by throwing pebbles? This question about the role and the value of statistics, in other contexts, is hotly debated. Fundamentally, we may suppose

Our first Monte Carlo simulation, on the Monte Carlo beach, generated 3156 hits for 4000 trials (see Table 1.1.3). We shall now see what this result tells us about  $\pi$ , at the most fundamental level of understanding. Hits and nonhits were generated by the Bernoulli distribution:

$$\xi_i = \begin{cases} 1 & \text{with probability } \theta \\ 0 & \text{with probability } (1 - \theta) \end{cases},$$
(24)

but the value of  $\pi/4 = \theta = \langle \xi_i \rangle$  is supposed unknown. Instead of the original variables  $\xi_i$ , we consider random variables  $\eta_i$  shifted by this unknown mean value:

$$\eta_i = \xi_i - \theta$$

The shifted random variables  $\eta_i$  now have zero mean and the same variance as the original variables  $\xi_i$ :

$$\langle \eta_i \rangle = 0$$
,  $\operatorname{Var} \eta_i = \operatorname{Var} \xi_i = \theta(1 - \theta) \le \frac{1}{4}$ .

Without invoking the limit  $N \to \infty$ , we can use the Chebyshev inequality to obtain an interval around zero containing at least 68% of the probability:

$$\left\{ \begin{array}{l} \text{with } 68\% \\ \text{probability} \end{array} \right\} : \quad \underbrace{\left| \frac{1}{N} \sum_{i=1}^{N} \eta_i \right|}_{\text{see eq. } (\ref{eq: see eq. } (\ref{eq: see$$

This has implications for the difference between our experimental result, 0.789, and the mathematical constant  $\pi$ . The difference between the two, with more than 68% chance, is smaller than 0.014:

$$\frac{\pi}{4} = 0.789 \pm 0.014 \Leftrightarrow \pi = 3.156 \pm 0.056, \tag{25}$$

where the value 0.056 is an upper bound for the 68% confidence interval that in physics is called an error bar. The quite abstract reasoning leading from eq. (24) to eq. (25)—in other words from the experimental result 3156 to the estimate of  $\pi$  with an error bar—is extremely powerful, and not always well understood. To derive the error bar, we did not use the central limit theorem, but the more general Chebyshev inequality. We also used an upper bound for the variance. With these precautions we arrived at the following result. We know with certainty that among an infinite number of beach parties, at which participants would play the same game of 4000 as we described in Sec. 1.1 and which would yield Monte Carlo results analogous to ours, more than 68% would hold the mathematical value of  $\pi$  inside their error bars. In arriving at this result, we did not treat the number  $\pi$  as a random variable—that would be nonsense, because  $\pi$  is a mathematical constant.

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