

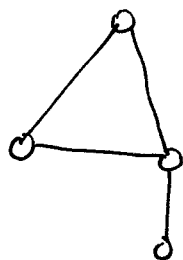
# Complexity barriers for partition functions

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$G$  bipartite  
 $\Downarrow$   
 $P(G, \mathbb{X}) = \text{per}(A_G)$

$$G = (V, E)$$

vertices edges



$e \in E \Rightarrow x_e$  variable associated with edge  $e$

associated with edge  $e$

$$\mathbb{X} = (x_e)_{e \in E}$$

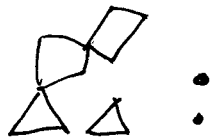
Ising partition function:

$$Z(G, \mathbb{X}) = \sum_{E' \subseteq E \text{ even}} \prod_{e \in E'} x_e$$

Dimer partition function:

$$P(G, \mathbb{X}) = \sum_{E' \subseteq E \text{ perfect matching}} \prod_{e \in E'} x_e$$

$E'$  even: each deg even



$E'$  perfect match.



substitution:

$$x_e := e^{-\beta J_e}$$

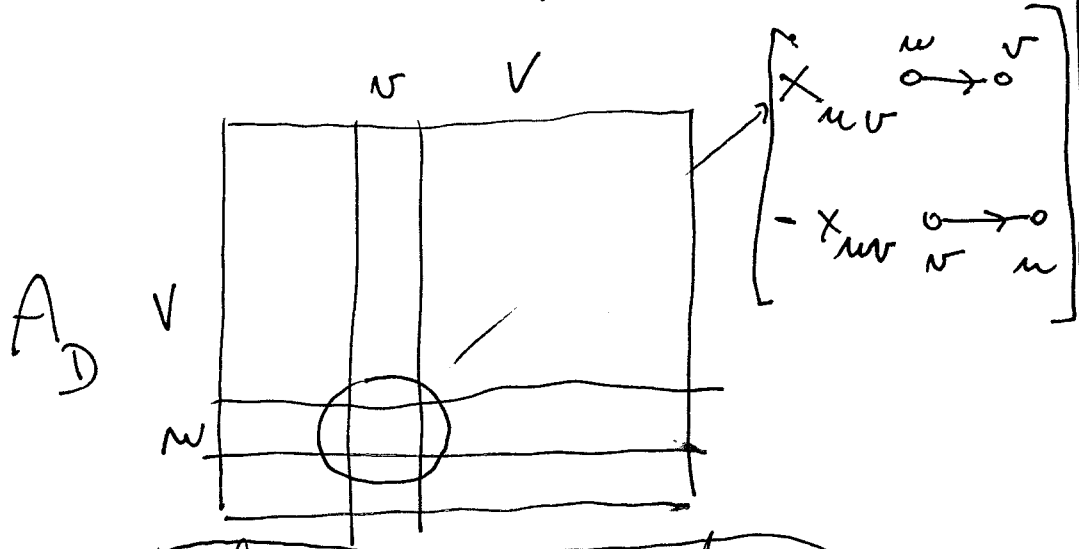
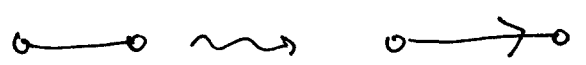
perfect matching

From 60's, basic tools are:

determinants  
(Pfaffian method)

Kasteleyn, Fisher, ...

$G \rightarrow D$  orientation



skew-symmetric

Pfaf  $A_D = (\text{Det } A_D)^{1/2}$

Products over aperiodic closed walks (Ihara-Selberg function)

discrete

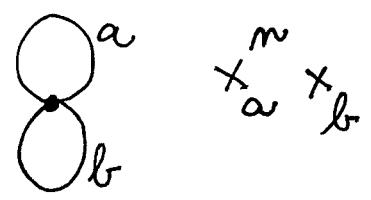
Kac, Ward, Feynman, Sherman ...

$\prod (1 - \prod x_e)$

rotation

$\prod$  aperiodic closed walks

infinite: formal power series



Bass' Theorem:  $\prod (1 - \prod x_e)$

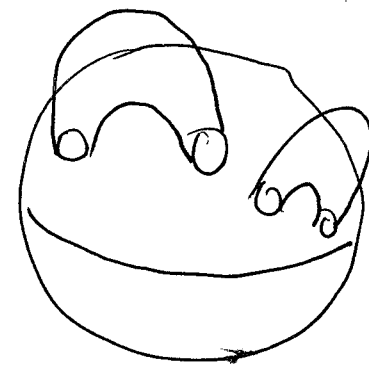
is a determinant

Complexity barriers of both tools : genus of  $G$

③

Theorem (Loeb, Masbaum, Adv. in Math 2011)

"Minimum number of orientations of  $G$  so that  $E(G, x)$  is linear combination of <sup>Pfaffians</sup> determinants of corresponding matrices is  $4^g$ ,  $g$  genus of  $G$ ."



$g=2$

- enumeration characterisation of genus
- exponential lower bound in very restricted computation model

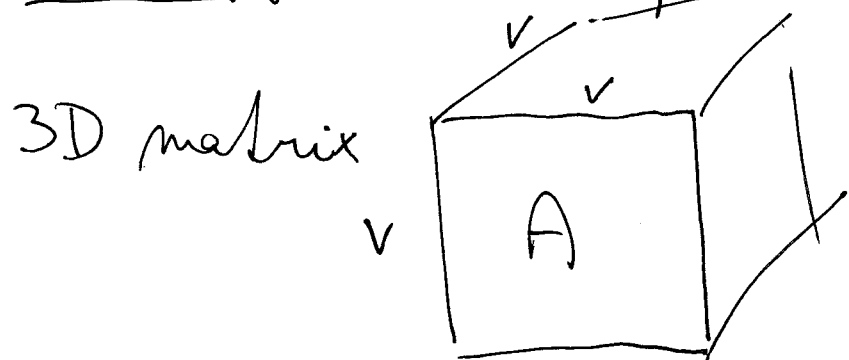
- still open for  $P(G, x)$  :  
 $4^g$  not true
- determinantal complexity: is ~~det~~ per exponentially harder than det (?)

How to deal with the genus barrier: study a "new" structure of the problems

\* Permanent has only combinatorial structure

\* (?) Polyhedral structure(?)

\* A suggestion: replace det by hyperdeterminants



$$\text{Per } A = \sum_{\pi_1, \pi_2} \prod_{i \in V} A_{i \pi_1(i) \pi_2(i)}$$

$$\text{Det } A = \sum_{\pi_1, \pi_2} \text{sign } \pi_1 \text{ sign } \pi_2 \prod_{i \in V} A_{i \pi_1(i) \pi_2(i)}$$

The same for 4D, 5D, ...

3D matrix  $A$  is Kasteleyn if there is signing  
 $A'$  of  $A$  so that  $\text{Per } A = \text{Det } A'$ .

\* (2D) ~~is~~ Kasteleyn matrices (rare): basically correspond  
 to planar graphs

\* Theorem (webl, Rykirk 2013)

- (3D) Kasteleyn matrices are (not rare).

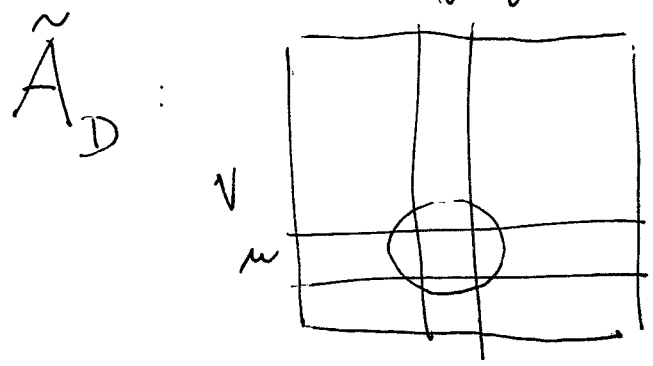
-  $G$  cubic lattice  $\Rightarrow$  there is 3D Kasteleyn matrix  $A_G$ :  
 $P(G, x) = \det A_G$ . Moreover  $A_G$  has "3D structure".

BUT: computing hyperdet is  
 #P complete (as hard as per)  
 few formulas

(?) Which advantage  
 (hyperdet) has over (per)  
 (?)

# Discrete Ihara-Selberg function

$G = (V, E)$ ,  $\mathcal{D}$  orientation of  $G$ ,



$x_e$  if  $u \xrightarrow{e} v$   
 $0$  otherwise

Theorem (Feynman, Sherman)

$$E^2(G, \mathbb{X}) = \prod_{\substack{\uparrow \text{ reduced} \\ \uparrow \text{ closed}}} \left[ 1 - (-1)^{\text{rot}} \prod_{e \in \mathcal{C}} x_e \right]$$

planar det (rot)

$$\otimes \det(I - \tilde{A}_{\mathcal{D}}) = \sum_{\mathcal{C} = \{c_1, \dots, c_\ell\} \text{ disjoint dicycles}} (-1)^\ell \prod_{e \in \mathcal{C}} x_e$$

$\otimes$  Theorem (Bass, Foata, Zeilberger ...)

$$\det(I - \tilde{A}_{\mathcal{D}}) = \prod_{\substack{\uparrow \text{ aperiodic} \\ \uparrow \text{ closed walk}}} \left( 1 - \prod_{e \in \mathcal{C}} x_e \right)$$

90's :

J. Distler, 3D Ising as a string theory

critical Ising :  
 discrete analogue of some conformal field theory [Alvarez-Gaumé, Moore, Vafa]  $(?)$   $g=2$   $(?)$

# 4D Bass' Theorem

Loebl 2014

(7)

$$(2D) \det(I - \tilde{A}_D)$$

incidence matrix  
of a digraph

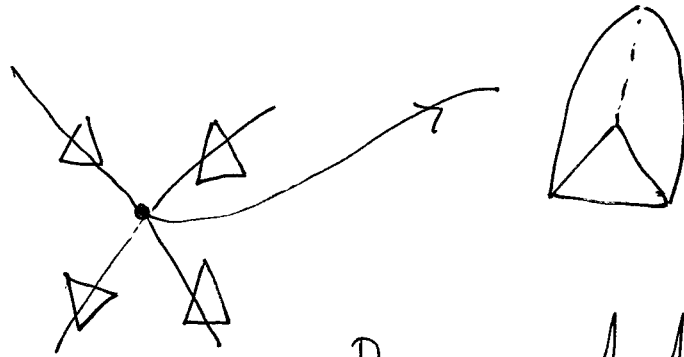
entries  $\longrightarrow$

graph: 1D simpl.  
complex

$$(4D) \det(I - A)$$

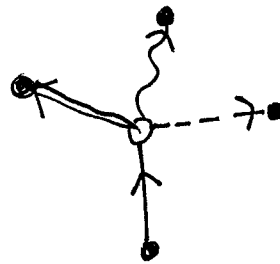
A: incidence matrix of orientation  
of 3D simplicial complex

entries directed tetrahedrons



orientation:  
each tetrahedron  
has 4 colored  $\Delta$

Representation:  $\circ$  vertex inside  
tetrahedron



$\bullet$  vertex inside  $\Delta$

# WALKS

(P)

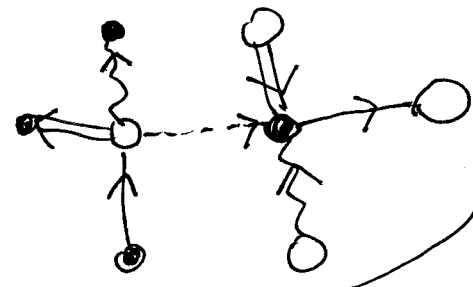
(2D)

closed connected  
aperiodic  
no beginning

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set of directed  
edges and  
"vertex connectors"  
with obvious  
properties

Branching and annihilation process on  $\Delta$   
through  $\Delta$ : closed connected aperiodic no  
beginning



set of directed tetrahedrons and  
 $\Delta$  connectors

FEASIBILITY PROPERTY: Process described from  
triangle  $t$ : at each time there is at most one  
active  $t$ -connector

- whose leaving arc is in the process
- whose entering arc is in the process

Active: at least 1 arc in process,  
- " - not in process.

